Three nucleon scattering in a "three dimensional" approach

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PWD VS 3D

- "CLASSICAL" APPROACH
- 3D APPROACH

2 GENERAL OPERATOR FORM

- ROTATIONS
- OTHER SYMMETRIES

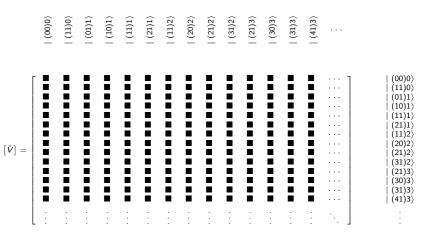
3N SCATTERING

CONVERGENCE

4 SUMMARY

PARTIAL WAVES - EXAMPLE

- For the moment let's focus on the 2N system.
- Let's try to calculate the 2N transition operator.



- \check{V} is the 2N potential.
- Each lives in a subspace with given orbital angular momentum *I*, spin *s* and total angular momentum *j* and different momentum magnitude states:

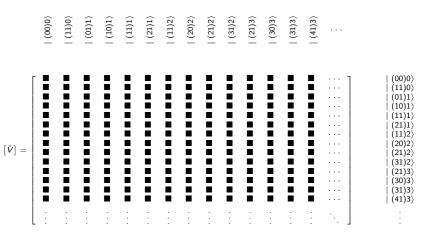
$$\langle |\mathbf{p}'|(l's')j'| \dots ||\mathbf{p}|(ls)j \rangle$$

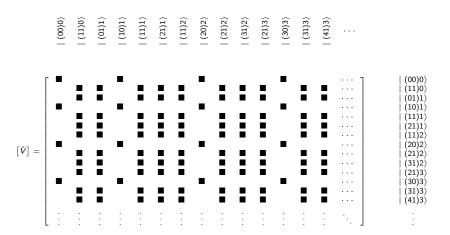
Impose pairity, time reversal and rotational symmetry ...

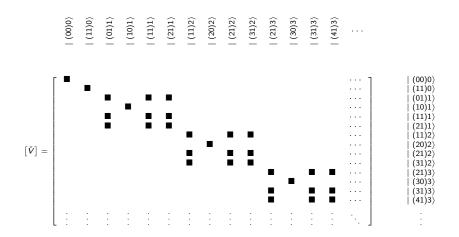
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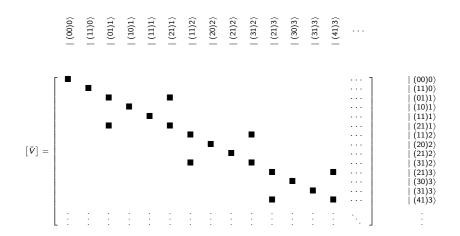
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Perform PWD on each operator of, eg., LSE t = V + V G₀t.
Solve the resulting linear equations.

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Perform PWD on each operator of, eg., LSE t = V + V G₀t.
Solve the resulting linear equations.

Battle tested.

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- Small numerical workload.
- Implementation requires heavily oscilating functions.
- It is not always obvious how many partial waves need to be taken into account.
- This is more complicated for three or more particles and different coupling schemes.
- Convergence problems for higher energies.

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- Lets take the 2N transition operator.
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- We would like to calculate the full transition operator. This is equivalent to calculating, for every p' and every p, the matrix element ⟨p' | ž | p⟩.
- This matrix element is an operator in spin space and has the form:

$$\left[\left< \mathbf{p}' \mid \check{t} \mid \mathbf{p} \right>
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$$\begin{bmatrix} t_{11}(p'_{x}, p'_{y}, p'_{z}, p_{x}, p_{y}, p_{z}) & t_{12}(\dots) & t_{13}(\dots) & t_{14}(\dots) \\ t_{21}(\dots) & t_{22}(\dots) & t_{23}(\dots) & t_{24}(\dots) \\ t_{31}(\dots) & t_{32}(\dots) & t_{33}(\dots) & t_{34}(\dots) \\ t_{41}(\dots) & t_{42}(\dots) & t_{43}(\dots) & t_{44}(\dots) \end{bmatrix}$$

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- We know that the solution has to satisfy appropriate symmetries.
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- The matrix element in momentum space can be written as a linear combination of 6 scalar functions t_i and spin operators [w_i]:

$$\left[\langle \mathbf{p}' \mid \check{t} \mid \mathbf{p} \rangle\right] = \sum_{i=1}^{6} t_i(|\mathbf{p}'|, |\mathbf{p}|, \hat{\mathbf{p}}' \cdot \hat{\mathbf{p}}) \left[\check{w}_i(\mathbf{p}, \mathbf{p})\right].$$

- Instead of calculating 16 functions of 6 real variables we now only need to calculate 6 functions of 3 variables.
- Couple orders of magnitude less numerical work!

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- Calculations can be easily modified to use different potentials.
- Operator fomrms (operators and states) significantly reduce numerical workload.
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If \check{R} is a spatial rotation, we require that operator \check{X} :

$$\check{R}^{-1}\check{X}\check{R}=\check{X}.$$

Ideally we would like to fit the operator into an operator form:

$$\langle \ldots \mid \check{X} \mid \ldots \rangle = \sum x_i \left[\check{O}_i\right].$$

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- Let's generalize a little bit and incorporate the dependance on the total momentum K.
- Boulding blocks (actually any number of momenta and spin vectors can be used):

 $\mathbb{T} = \{\check{\mathbf{p}}', \check{\mathbf{p}}, \check{\mathbf{K}}, \check{\boldsymbol{\sigma}}(1), \check{\boldsymbol{\sigma}}(2)\}.$

In principle we have to consider all scalar combinations of the elements from T. For example if v_i ∈ T we could use:

$$(\check{\boldsymbol{v}}_1\times(\check{\boldsymbol{v}}_2\times\check{\boldsymbol{v}}_3))\cdot(\check{\boldsymbol{v}}_4\times(\check{\boldsymbol{v}}_5\times(\check{\boldsymbol{v}}_6\times\check{\boldsymbol{v}}_7)))$$

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Simple vector identities lead to:

$$\begin{split} (\check{\mathbf{v}}_1 \times (\check{\mathbf{v}}_2 \times \check{\mathbf{v}}_3)) \cdot (\check{\mathbf{v}}_4 \times (\check{\mathbf{v}}_5 \times (\check{\mathbf{v}}_6 \times \check{\mathbf{v}}_7))) \\ &= (\check{\mathbf{v}}_1 \cdot \check{\mathbf{v}}_3) (\check{\mathbf{v}}_2 \cdot \check{\mathbf{v}}_5) (\check{\mathbf{v}}_4 \times \check{\mathbf{v}}_6 \cdot \check{\mathbf{v}}_7) \\ &- (\check{\mathbf{v}}_1 \cdot \check{\mathbf{v}}_2) (\check{\mathbf{v}}_3 \cdot \check{\mathbf{v}}_5) (\check{\mathbf{v}}_4 \times \check{\mathbf{v}}_6 \cdot \check{\mathbf{v}}_7) \\ &+ (\check{\mathbf{v}}_1 \cdot \check{\mathbf{v}}_2) (\check{\mathbf{v}}_4 \cdot \check{\mathbf{v}}_5) (\check{\mathbf{v}}_3 \times \check{\mathbf{v}}_6 \cdot \check{\mathbf{v}}_7) \\ &- (\check{\mathbf{v}}_1 \cdot \check{\mathbf{v}}_3) (\check{\mathbf{v}}_4 \cdot \check{\mathbf{v}}_5) (\check{\mathbf{v}}_2 \times \check{\mathbf{v}}_6 \cdot \check{\mathbf{v}}_7) \end{split}$$

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■ A general observation can be made: Any scalar expression constructed from operators in T can be constructed from a combination of operators in the set V:

$$\mathbb{V} = \{\check{1}, \check{\mathbf{v}}_i \cdot \check{\mathbf{v}}_j , (\check{\mathbf{v}}_i \times \check{\mathbf{v}}_j) \cdot \check{\mathbf{v}}_k \}.$$

For example, from the previous slide, we have the following CHAINS of operators of length 3:

$$\begin{split} \begin{bmatrix} \tilde{\zeta}_1 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_5) (\tilde{v}_4 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_2 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_2) (\tilde{v}_3 \cdot \tilde{v}_5) (\tilde{v}_4 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_3 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_2) (\tilde{v}_4 \cdot \tilde{v}_5) (\tilde{v}_3 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_4 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_4 \cdot \tilde{v}_5) (\tilde{v}_2 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_5 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_6) (\tilde{v}_2 \cdot \tilde{v}_3) (\tilde{v}_4 \cdot \tilde{v}_5), \\ \begin{bmatrix} \tilde{\zeta}_6 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_6) (\tilde{v}_4 \cdot \tilde{v}_5), \\ \begin{bmatrix} \tilde{\zeta}_7 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_5) (\tilde{v}_4 \cdot \tilde{v}_6), \\ \begin{bmatrix} \tilde{\zeta}_8 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_5) (\tilde{v}_4 \cdot \tilde{v}_6). \end{split}$$

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$$\begin{split} \begin{bmatrix} \tilde{\zeta}_1 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_5) (\tilde{v}_4 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_2 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_2) (\tilde{v}_3 \cdot \tilde{v}_5) (\tilde{v}_4 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_3 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_2) (\tilde{v}_4 \cdot \tilde{v}_5) (\tilde{v}_3 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_4 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_4 \cdot \tilde{v}_5) (\tilde{v}_2 \times \tilde{v}_6 \cdot \tilde{v}_7), \\ \begin{bmatrix} \tilde{\zeta}_5 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_6) (\tilde{v}_2 \cdot \tilde{v}_3) (\tilde{v}_4 \cdot \tilde{v}_5), \\ \begin{bmatrix} \tilde{\zeta}_6 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_6) (\tilde{v}_4 \cdot \tilde{v}_5), \\ \begin{bmatrix} \tilde{\zeta}_7 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_5) (\tilde{v}_4 \cdot \tilde{v}_6), \\ \begin{bmatrix} \tilde{\zeta}_8 \end{bmatrix} &= (\tilde{v}_1 \cdot \tilde{v}_3) (\tilde{v}_2 \cdot \tilde{v}_5) (\tilde{v}_4 \cdot \tilde{v}_6). \end{split}$$

■ A general observation can be made: Any scalar expression constructed from operators in T can be constructed from a combination of operators in the set V:

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- How get from this to the operator form?

$$\langle \mathbf{p}'\mathbf{K} \mid \check{X} \mid \mathbf{p}\mathbf{K} \rangle = \ldots + \sum_{i=1}^{8} x_i \left[\check{C}_i\right] + \ldots$$

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- Start with all chains of length 1.
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After the first iteration 11 unique chains of length 1:

$$\begin{split} \tilde{I} \\ \tilde{p}' \cdot \tilde{\sigma}(1) \\ \tilde{p}' \cdot \tilde{\sigma}(2) \\ \tilde{p} \cdot \tilde{\sigma}(1) \\ \tilde{p} \cdot \tilde{\sigma}(2) \\ \tilde{K} \cdot \tilde{\sigma}(1) \\ \tilde{K} \cdot \tilde{\sigma}(2) \\ \tilde{\sigma}(1) \cdot \tilde{\sigma}(2) \\ \tilde{\sigma}'(1) \cdot \tilde{\sigma}(2) \\ \tilde{p} \times \tilde{\sigma}(1)) \cdot \tilde{\sigma}(2) \\ \tilde{\kappa} \times \tilde{\sigma}(1)) \cdot \tilde{\sigma}(2). \end{split}$$

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After the second iteration 16 independent chains of length < 2. Since this is the last iteration, they have special names.

 $\left[\check{O}_{6}\right] = \check{\mathsf{K}} \cdot \check{\sigma}(1)$ $[\check{O}_7] = \check{K} \cdot \check{\sigma}(2)$

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The third iteration does not introduce any new unique operators!
We just created the operator form for X:

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- Let D be a grup of transformations constructed from **parity**, **time reversal**, **Hermitian conjugate and two particle exchange**.
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N-d elastic scattering and breakup description via the 3N Faddeev equation

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Calculate observables in the breakup channel ($\langle \phi_o |$ - three free nucleons, $| \phi \rangle$ - deuteron and free nucleon):

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Compare 3D approach and "battle tested" PWD approach.

ELASTIC SCATTERING:

BREAKUP:





- Initially the neutron: q_i . In the final state, the neutron: $q_{f^{(i)}}$
- Scattering parametrized by θ_{c.m.}.

- In the final state the Jacobi momenta: p^t and q^t.
- Scattering parametrized by the kinematic curve parameter *S*: **p**^{*f*}(*S*), **q**^{*f*}(*S*).

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ELASTIC SCATTERING:







- **I** Initially the neutron: \mathbf{q}_i . In the final state, the neutron: \mathbf{q}_f .
- Scattering parametrized by θ_{c.m.}.

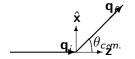
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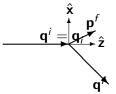
^eFB²³ AARHUS

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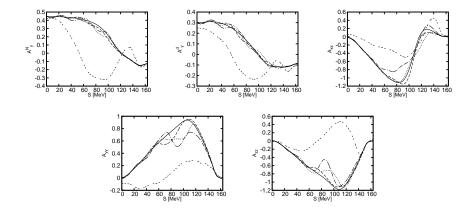
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- Results for breakup [Eur. Phys. J. A 51:132 (2015)].
- Deuteron and nucleon vector analyzing powers (A^d_y, A^N_y) and the deuteron tensor analyzing powers (A_{xx}, A_{yy}, A_{zz}) LAB energy 190 MeV.
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- The dashed-dotted, dotted and dashed lines PWD results with max. total anguler momentum 21/2, 23/2, 25/2 and max. 2 – 3 angular momentum 8.

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- First order results for neutron deuteron scattering suggest that the 3D approach can be used to achieve convergence at higher energies.
- There is a necessity to construct new general operator forms.
- Constructing $\check{T} \mid \phi \rangle$ can lead to efficient calculations.
- Possibility to add relativistic corrections to the calculations.

THANK YOU

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